

## Appendum to Identification of Conditional Interventional Distributions

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function **c-identify**( $C, T, Q[T]$ )  
INPUT:  $T, C \subseteq T$  are both are C-components,  
 $Q[T]$  a probability distribution  
OUTPUT: Expression for  $Q[C]$  in terms of  $Q[T]$  or **FAIL**

let  $A = An(C)_{G_T}$

- 1 if  $A = C$ , return  $\sum_{T \setminus C} P$
- 2 if  $A = T$ , return **FAIL**
- 3 if  $C \subset A \subset T$ , there exists a C-component  $T'$  such that  $C \subset T' \subset A$ . return **c-identify**( $C, T', Q[T']$ )  
( $Q[T']$  is known to be computable from  $\sum_{T \setminus A} Q[T]$ )

Figure 1: A C-component identification algorithm from [Tian, 2004].

The following algorithm, **cond-identify**, appears in [Tian, 2004]. A proof in [Shpitser & Pearl, 2006a] claims this algorithm is not sound, but was based on a misunderstanding of notation. In fact, this algorithm can be shown to be complete.

**Theorem 1** *cond-identify is complete.*

*Proof:* We want to show that whenever the algorithm fails, the corresponding effect is not identifiable. The operation of the algorithm can be thought of as follows – it starts with a set of 'bad' C-components (containing a non-identifiable effect). These 'bad' C-components 'infect' other C-components which were initially 'good' (identifiable). The 'infection' proceeds until no more C-components can be 'infected' or until we encounter the set  $\mathbf{Y}$  of effect variables. In the latter case the algorithm fails.

Our proof is by induction.

A known result in [Shpitser & Pearl, 2006b] is that if **identify** fails to identify  $D_i$  from  $C_i$  then  $D_i$  is not identifiable. If the algorithm fails, such a  $D_i$  is guaranteed to exist, and is in the ancestor set of  $\mathbf{Y} \cup \mathbf{Z}$ . Fix such a  $D_i$ .

function **cond-identify**( $\mathbf{y}, \mathbf{x}, \mathbf{z}, P, G$ )

INPUT:  $\mathbf{x}, \mathbf{y}, \mathbf{z}$  value assignments,  $P$  a probability distribution,  $G$  a causal diagram (an I-map of  $P$ ).

OUTPUT: Expression for  $P_{\mathbf{x}}(\mathbf{y}|\mathbf{z})$  in terms of  $P$  or **FAIL**.

- 1 let  $D = An(\mathbf{Y} \cup \mathbf{Z})_{G_{\mathbf{X}}}, F = D \setminus (\mathbf{Y} \cup \mathbf{Z})$
- 2 assume  $C(D) = \{D_1, \dots, D_k\}$
- 3 let  $N = \{D_i | \mathbf{c-identify}(D_i, C_{D_i}, Q[C_{D_i}]) = \mathbf{FAIL}\}$
- 4 if  $N = \emptyset$ , return  $\frac{\sum_f \prod_i Q[D_i]}{\sum_{\mathbf{y}, \mathbf{f}} \prod_i Q[D_i]}$
- 5 let  $F_0 = F \cap (\bigcup_{D_i \in N} Pa(D_i)), I = C(D) \setminus N$
- 6 remove the set  $\{D_i | Pa(D_i) \cap F_0 \neq \emptyset\}$  from  $I$  and add it to  $I_0$  (which is initially empty)
- 7 let  $B = (F \setminus F_0) \cap \bigcup_{D_i \in I_0} Pa(D_i)$
- 8 if  $B \neq \emptyset$ , add all nodes in  $B$  to  $F_0$ , and go to line 6
- 9 if  $\mathbf{Y} \cap (\bigcup_{D_i \in (N \cup I_0)} Pa(D_i)) \neq \emptyset$ , return **FAIL**,  
else return  $\frac{\sum_{\mathbf{f}_1} \prod_{D_i \in I_1} Q[D_i]}{\sum_{\mathbf{y}, \mathbf{f}_1} \prod_{D_i \in I_1} Q[D_i]}$

Figure 2: An identification algorithm from [Tian, 2004]. For each  $D_i$ , we denote  $C_{D_i} \in C(G)$  such that  $D_i \subseteq C_{D_i}$ .

Fix the minimal set  $\mathbf{W} \subseteq \mathbf{Y} \cup \mathbf{Z}$  such that  $D_i \in An(\mathbf{W})_{G_{\mathbf{Y}, \mathbf{Z}}}$ . Our base case will be that  $\mathbf{Y} \cap \mathbf{W} \neq \emptyset$ . Note that because  $D_i$  is a C-component, every element in  $\mathbf{Z} \cap \mathbf{W}$  has a backdoor path to  $\mathbf{Y} \cap \mathbf{W}$ . Then our conclusion follows by the backdoor hedge criterion [Shpitser & Pearl, 2006a].

Assume  $\mathbf{Y} \cap \mathbf{W} = \emptyset$ . Since the algorithm failed,  $\mathbf{Y}$  intersects the set of 'infected' C-components. We want to show that  $\mathbf{W}$  has a backdoor path to  $\mathbf{Y}$ , which follows the 'infected' C-components. Our conclusion will then follow by the backdoor hedge criterion.

Without loss of generality, assume all variables outside  $C_i$  are observable. That means all C-components that are not  $C_i$  contain one variable. We prove inductively that all 'infected' nodes are d-connected. The base case is nodes in  $D_i$ . Note that each node in  $D_i$  has a descendant in  $\mathbf{Z}$ . Since  $D_i$  is a C-component, each pair of nodes in  $D_i$  are d-connected by a bidirected path.

For the inductive hypothesis, consider a set of infected nodes  $N$  which are in the ancestor set of  $\mathbf{Z}$  (but not  $\mathbf{Y}$ ) and which are pairwise d-connected. A new node  $I$  can become 'infected' in one of three ways:

If  $I$  and  $N$  share a parent which is not in  $\mathbf{Z} \cup \mathbf{Y}$ , that means some node in  $N$  has a d-connected path to  $I$  through the common parent. Since this node in  $N$  is an ancestor of  $\mathbf{Z}$ , this path can be extended to a d-connected path to any node in  $N$ . If  $I$  is a parent of some node in  $N$ , that node has a d-connected path to  $I$ . Moreover, since that node is an ancestor of  $\mathbf{Z}$ , that path can be extended to a d-connected path to any node in  $N$ . If some node in  $N$  is a parent of  $I$ , the reasoning is the same.

$I$  itself can either be in  $\mathbf{Y}$ , an ancestor of  $\mathbf{Y}$ , or an ancestor of  $\mathbf{Z}$ . In the first case, we are done since we constructed a d-connected path from a parent of  $\mathbf{Z}$  to a node in  $\mathbf{Y}$ , which translates into a backdoor path from  $\mathbf{Z}$  to  $\mathbf{Y}$ . In the second case, the d-connected path from a parent of  $\mathbf{Z}$  to an ancestor of  $\mathbf{Y}$  easily extends to a backdoor path from  $\mathbf{Z}$  to  $\mathbf{Y}$ . In the last case, we simply continue the induction until we reach either of the first two cases. We know we reach these cases eventually since the algorithm failed.  $\square$

## References

- [Shpitser & Pearl, 2006a] Shpitser, I., and Pearl, J. 2006a. Identification of conditional interventional distributions. In *Uncertainty in Artificial Intelligence*, volume 22.
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